Key Management and Distribution Cloud-Computing Security

Ruben Niederhagen

May 10, 2013

"Key management is the hardest part of cryptography and often the Achilles' heel of an otherwise secure system."

Bruce Schneier

1. Terminology

Symmetric Key Distribution
 Using Symmetric Encryption
 Using Asymmetric Encryption

3. Public Key Distribution
X.509 Certificates
Public Key Infrastructure
Risks

4. Key Management Strategies in the Cloud

Core Principles of Information Security

- Confidentiality
- Integrity
- Availability

Parkerian Hexad by Donn B. Parker in 2002

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- Integrity
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- Authenticity
 - Non-Repudiation
 - Plausible Deniability

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 - Plausible Deniability

Anonymity, (Perfect) Forward Secrecy, Trust, ...

Building Blocks

Symmetric Encryption

Using the same shared key for encryption and decryption:

$$C = E(K, P)$$
 $P = E(K, C)$

Examples: Twofish, Serpent, AES (Rijndael), Blowfish, 3DES, ...

Building Blocks

Symmetric Encryption

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Examples: Twofish, Serpent, AES (Rijndael), Blowfish, 3DES, ...

Asymmetric Encryption

Using public key for encryption and private key for decryption:

$$C = E(K_{pub}, P)$$
 $P = E(K_{priv}, C)$

Examples: RSA, McEliece, ElGamal, ...

Why do we need key distribution?



Secure communication requires shared, "a priori" knowledge.

How do we achieve this knowledge?

By using some kind of key-distribution scheme!

Key Distribution Using Symmetric Encryption

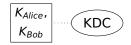
Given parties A and B, there are several alternatives for key distribution

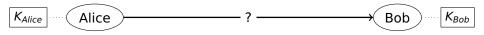
- 1. A can select key and physically deliver it to B.
- 2. A third party can select and physically deliver the key to A and B.
- 3. If A and B have communicated previously, they can use the previous key to encrypt a new key.
- 4. If A and B have secure communication channel with a third party C, C can relay the key between A and B.

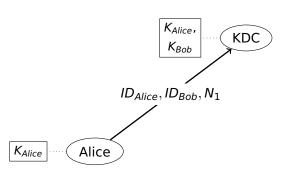
Key Distribution Using Symmetric Encryption

Key Hierarchy

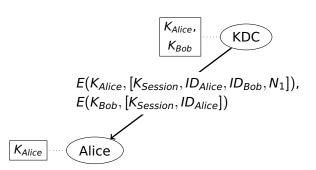
- typically have a hierarchy of keys
- session key
 - temporary key
 - used for encryption of data between users for one logical session
 - discarded after usage
- master key
 - longterm key
 - used to encrypt session keys
 - shared by user and key distribution center



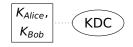


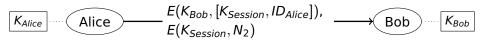


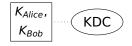




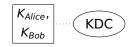




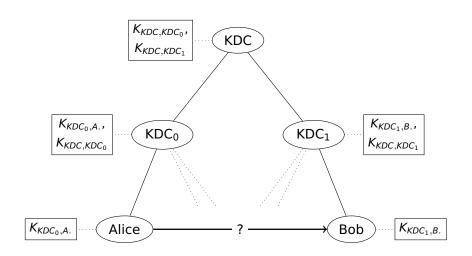


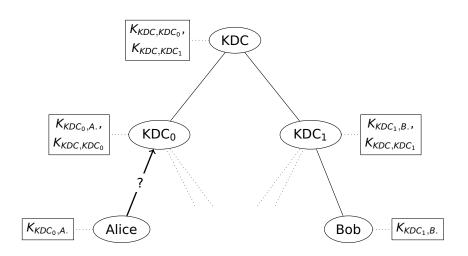


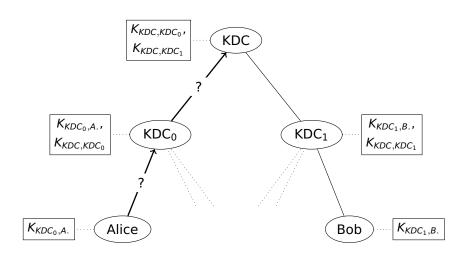


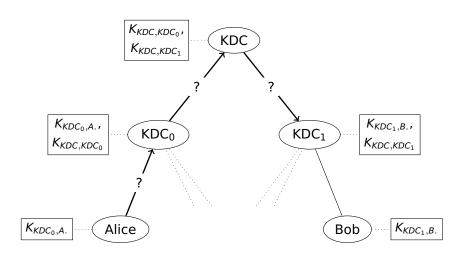


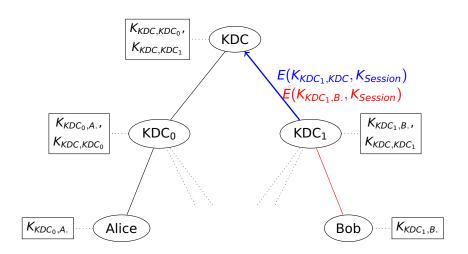


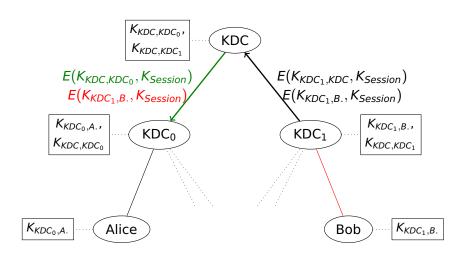


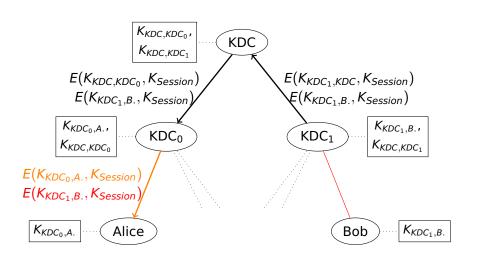


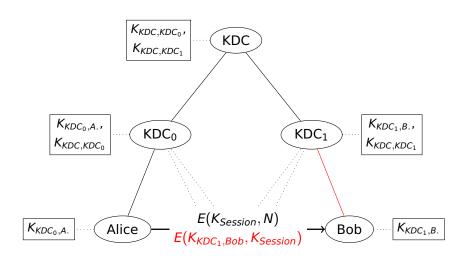


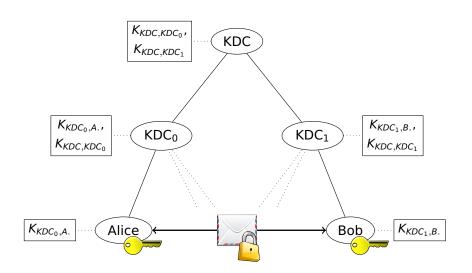












Session-Key Lifetime

Make Lifetime as Short as Possible

- benefit:
 - reduced attack surface
 - less information compromised in case encryption is broken
- disadvantage:
 - requires to obtain keys more often
 - requires more time

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Connection-Oriented Protocols

- naturally choice: one key per connection
- re-key if connection is maintained for too long

Session-Key Lifetime

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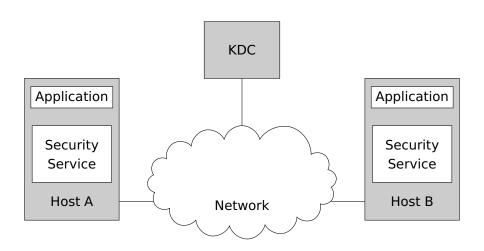
- benefit:
 - reduced attack surface
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- disadvantage:
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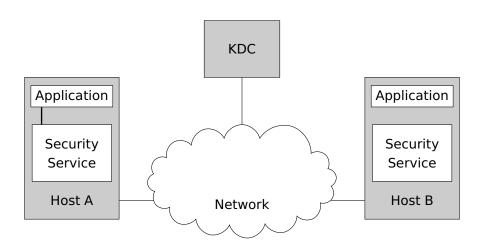
Connection-Oriented Protocols

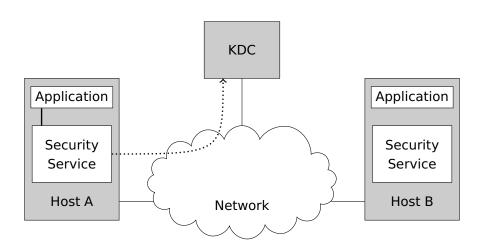
- naturally choice: one key per connection
- re-key if connection is maintained for too long

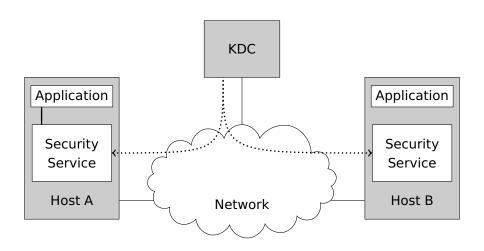
Connectionless Protocols

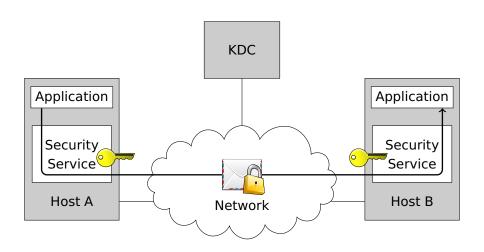
- naturally choice: one key per exchange
- re-key after a certain amount of time

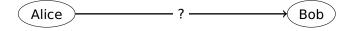


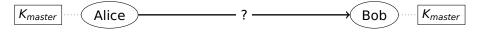


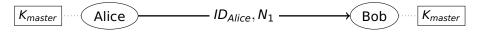


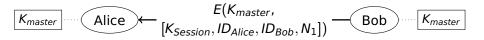












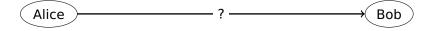


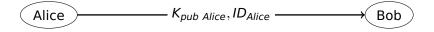
Requires n(n-1)/2 master keys for n nodes.

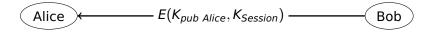
Key Distribution Using Asymmetric Encryption

Features of Asymmetric Encryption Schemes

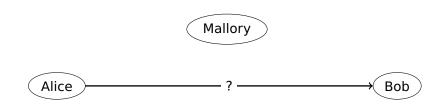
- typically slower than symmetric schemes
- can be used to encrypt symmetric keys for distribution
- public key can be distributed openly

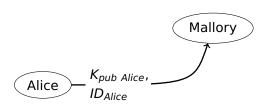






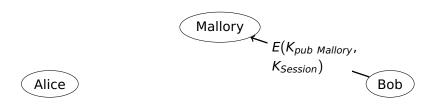






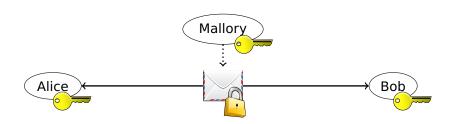
Bob

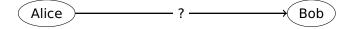


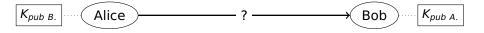


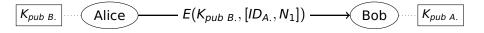


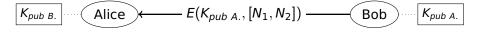
Bob

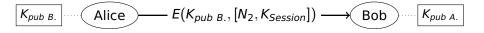








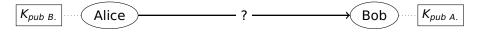








Use signatures!



$$\begin{array}{ccc}
E(K_{pub\ Bob}, K_{Session}), \\
N, ID_{Alice}, \\
E(K_{priv\ Alice}, f([...]))
\end{array}$$
Bob $K_{pub\ A}$



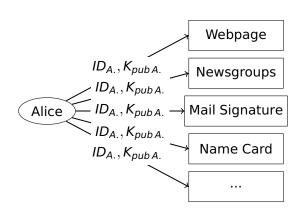


How do we distribute the public keys?

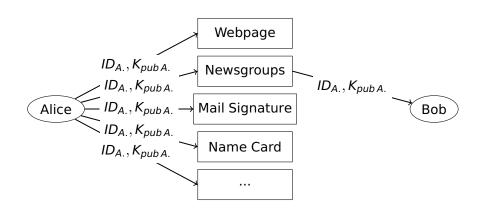
Public Key Distribution

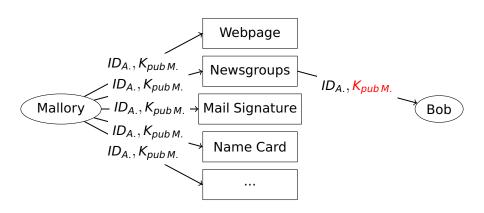
General Schemes:

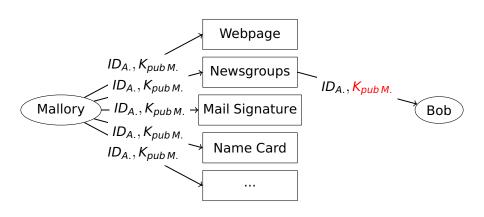
- public announcement
- publicly available directory
- public-key authority
- public-key certificates



Bob







Easy to forge!

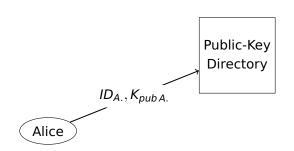
Public-Key Directory

Alice

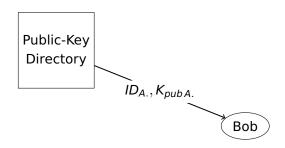
Bob

Public Key Distribution

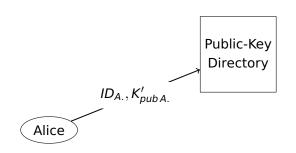
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Bob

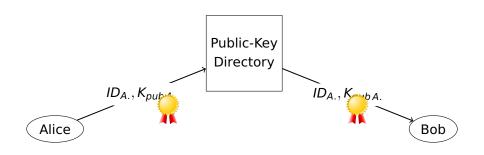






Bob

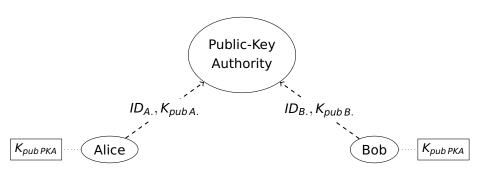
Publicly Available Directory



Communication with public-key directory must be authenticated, acknowledged, and protected against replay attacks!

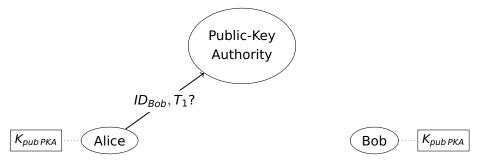
Public Key Distribution May 10, 2013

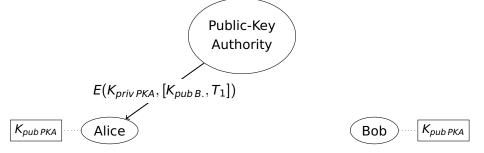
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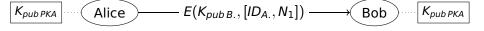


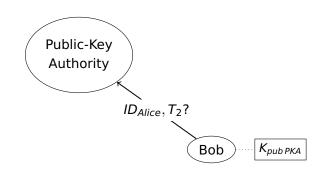
K_{pub PKA} Alice ? Bob K_{pub PKA}



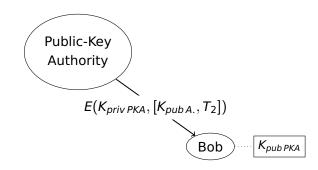
















$$K_{pub\,PKA}$$
 Alice $E(K_{pub\,A.},[N_1,N_2])$ Bob $K_{pub\,PKA}$



$$K_{pub\ PKA}$$
 Alice $E(K_{pub\ B.}, [N_2, K_{Session}])$ Bob $K_{pub\ PKA}$





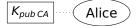




Long latencies due to communication with PKA!

- Alice and Bob may cache public keys.
- Use certificates...

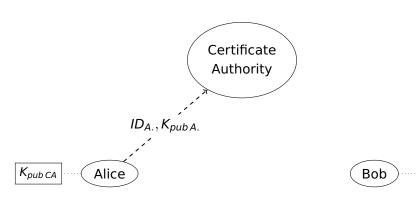




Bob K_{pub CA}

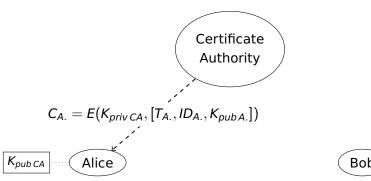
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Public Key Distribution May 10, 2013

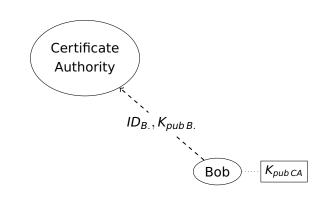


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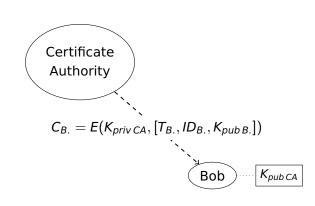
 $K_{pub CA}$



Bob Kpub CA





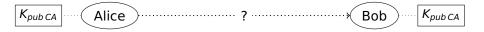


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Public Key Distribution May 10, 2013







$$K_{pub CA}$$
 Alice $C_{A.} = E(K_{priv CA}, [T_{A.}, ID_{A.}, K_{pub A.}])$ Bob $K_{pub CA}$



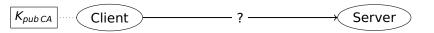




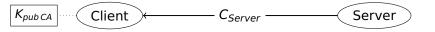














$$Client \longrightarrow E(K_{pub \, Server}, K_{Session}) \longrightarrow Server$$

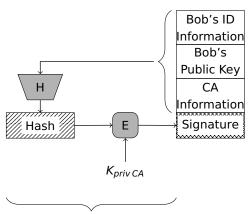




X.509 Certificates

Bob's ID Information Bob's Public Key CA Information

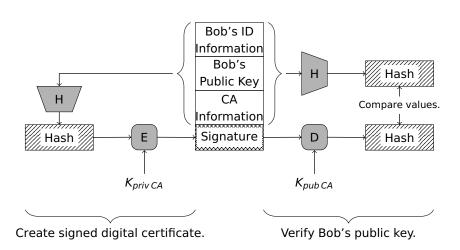
X.509 Certificates



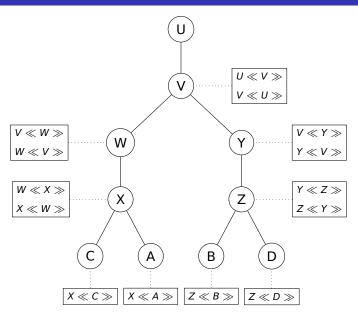
Create signed digital certificate.

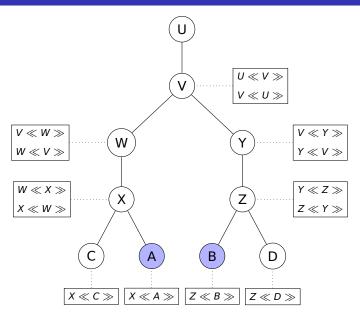
Public Key Distribution X.509 Certificates May 10, 2013 24 / 34

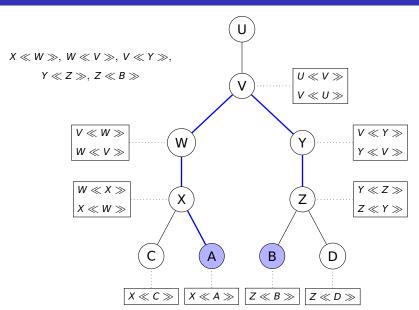
X.509 Certificates



Public Key Distribution X.509 Certificates May 10, 2013 24 / 34

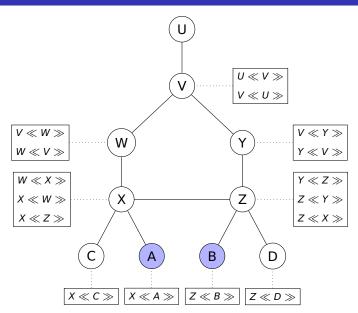


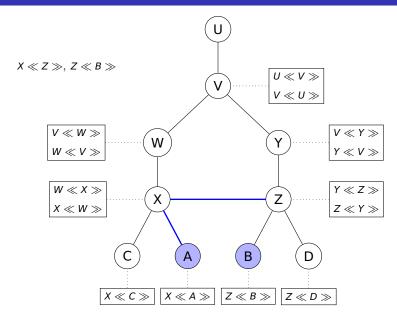




Public Key Distribution X.509 Certificates May 10, 2013

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Public Key Distribution X.509 Certificates May 10, 2013

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X.509 Example

```
Certificate:
  Data:
       Version: 1 (0x0)
                                                     Server Certificate
       Serial Number: 7829 (0x1e95)
       Signature Algorithm: md5WithRSAEncryption
       Issuer: C=ZA, ST=Western Cape, L=Cape Town, ...
       Validity
           Not Before: Jul 9 16:04:02 1998 GMT
           Not After: Jul 9 16:04:02 1999 GMT
       Subject: C=US, ST=Maryland, L=Pasadena, O=Brent Baccala, ...
       Subject Public Key Info:
           Public Key Algorithm: rsaEncryption
           RSA Public Key: (1024 bit)
               Modulus (1024 bit):
                   00:b4:31:98:0a:c4:bc:62:c1:88:aa:dc:b0:c8:bb:...
               Exponent: 65537 (0x10001)
  Signature Algorithm: md5WithRSAEncryption
       93.5f.8f.5f.c5.af.hf.0a.ah.a5.6d.fh.24.5f.h6.59.5d.9d.
```

Public Key Distribution X.509 Certificates May 10, 2013 26 / 34

X.509 Example

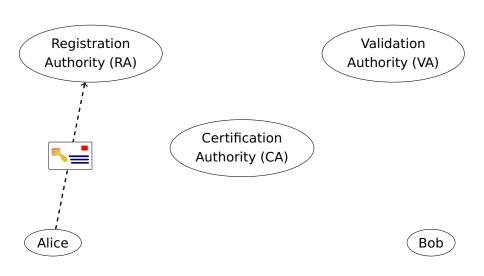
```
Certificate:
   Data:
       Version: 3 (0x2)
                                                     Sub-CA Certificate
       Serial Number: 1 (0x1)
       Signature Algorithm: md5WithRSAEncryption
       Issuer: C=ZA, ST=Western Cape, L=Cape Town, ...
       Validity
           Not Before: Jul 9 16:04:02 1998 GMT
           Not After: Jul 9 16:04:02 1999 GMT
       Subject: C=US, ST=Maryland, L=Pasadena, O=Brent Baccala, ...
       Subject Public Key Info:
           Public Key Algorithm: rsaEncryption
           RSA Public Key: (1024 bit)
               Modulus (1024 bit):
                   00:b4:31:98:0a:c4:bc:62:c1:88:aa:dc:b0:c8:bb:...
               Exponent: 65537 (0x10001)
       X509v3 extensions:
           X509v3 Basic Constraints: critical
               CA: TRUE
   Signature Algorithm: md5WithRSAEncryption
       93:5f:8f:5f:c5:af:hf:0a:ah:a5:6d:fh:24:5f:h6:59:5d:9d:...
```

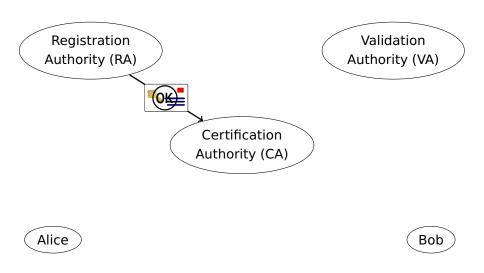
Registration Authority (RA) Validation Authority (VA)

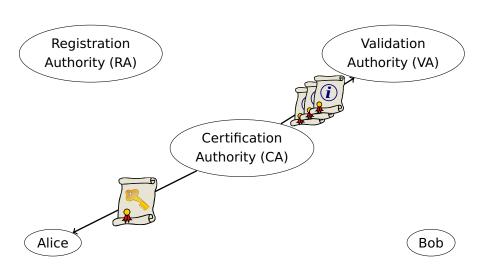
Certification
Authority (CA)

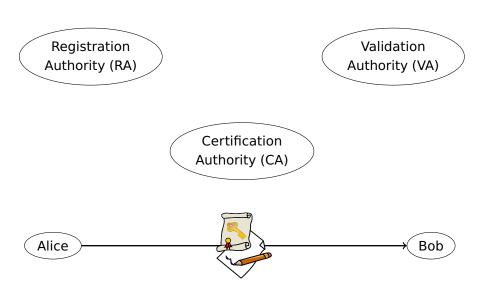
Alice

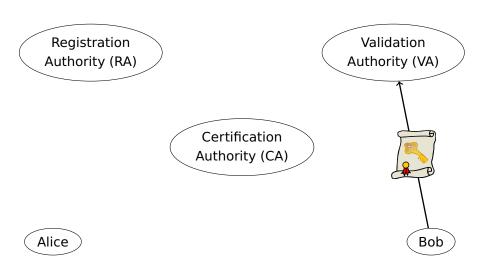
Bob

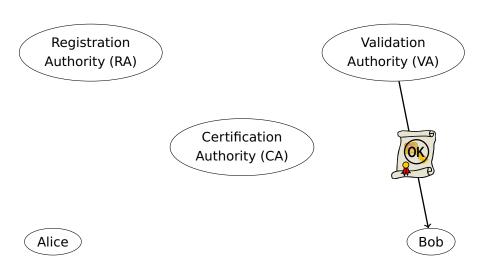


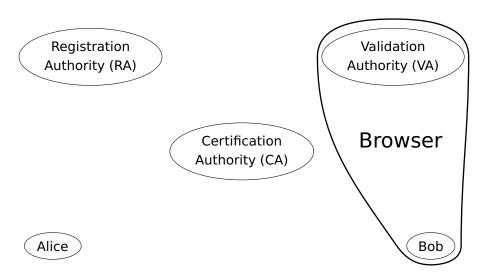












Comodo Security Breach, March 15, 2011

- A user account with an affiliate registration authority had been compromised.
- The attacker issued nine certificate signing requests.
- Certificates for issued for:
 - mail.google.com
 - login.live.com
 - www.google.com
 - login.yahoo.com
 - login.skype.com
 - addons.mozilla.org
- ► The attack was traced to IP address 212.95.136.18 in Tehran, Iran.
- The origin of the attack may be the "result of an attacker attempting to lay a false trail."

Public Key Distribution Risks May 10, 2013 28 / 34

DigiNotar Fraudulent Certificates, July 10, 2011

- An attacker hacked into the systems of DigiNotar and issued a certificate for Goolge.
- ► This certificate was subsequently used by unknown persons in Iran to conduct a man-in-the-middle attack against Google services.
- After this certificate was found, DigiNotar belatedly admitted dozens of fraudulent certificates had been created, including certificates for the domains of Yahoo!, Mozilla, WordPress and The Tor Project.
- ► Google blacklisted 247 certificates in Chromium, but the final known total of misissued certificates is at least 531.
- DigiNotar also controlled an intermediate certificate which was used for issuing certificates as part of the Dutch government's public key infrastructure "PKloverheid" program.

TURKTRUST and the Chain of Trust, 2012

- ► TURKTRUST sent two intermediate certificates to organisations that had requested regular certificates.
- One of the certificates was revoked at the request of the customer who received it.
- ► The other organisation now had the ability to sign SSL certificates for any domain name it chose.
- There is no known malicious use of this intermediate certificate
 but TURKTRUST should never have issued it in the first place.

► Comodo:

DigiNotar:

► TURKTRUST:

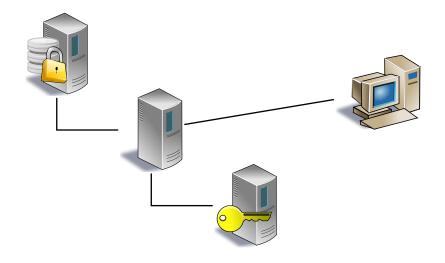
- ▶ Comodo:
 - Authenticity and legitimacy of signing request.
- DigiNotar:
- ► TURKTRUST:

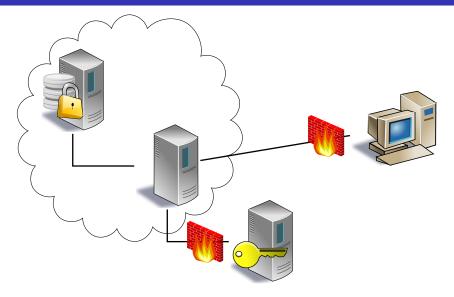
Public Key Distribution Risks May 10, 2013 31 / 34

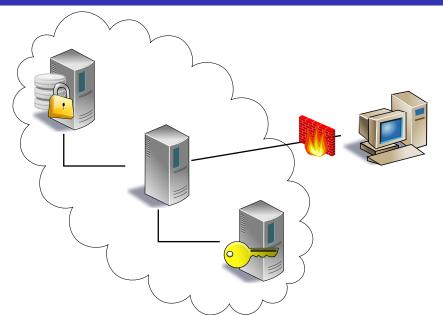
- Comodo:
 - Authenticity and legitimacy of signing request.
- DigiNotar:
 - Security and integrity of certificate authority.
- ► TURKTRUST:

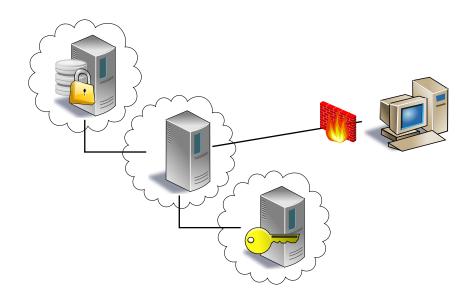
- Comodo:
 - Authenticity and legitimacy of signing request.
- DigiNotar:
 - Security and integrity of certificate authority.
- ► TURKTRUST:
 - Just do it right.

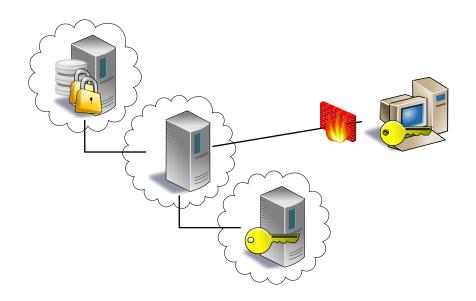
Public Key Distribution Risks May 10, 2013 31 / 34

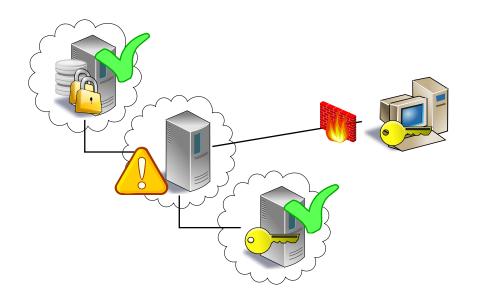












Fully Homomorphic Encryption

- May be used to protect data during computation in the Cloud.
- ► Far away from being practical; might never be feasible.

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- ► Far away from being practical; might never be feasible.

"Visions of a fully homomorphic cryptosystem have been dancing in cryptographers' heads for thirty years. I never expected to see one. It will be years before a sufficient number of cryptographers examine the algorithm that we can have any confidence that the scheme is secure."

Bruce Schneier

References

Further Reading:

William Stallings, *Cryptography and Network Security: Principles and Practice*, 5th ed. Prentice Hall, Upper Saddle River, NJ, USA. January, 2010.

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